

Polynomials in homotopy type theory

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LHC days 2024

The situation:

- Polynomials in a category are a categorification of ordinary polynomials

$$F(X) = X \times X + 1$$

- They can be defined in any locally cartesian closed category
- Similar to combinatorial species and their generalizations

Our contribution:

- Show that polynomials are Kleisli morphisms for a comonad on spans
- and fit in a model of linear logic
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- Decomposing $(A \Longrightarrow B)$ as $(!A \multimap B)$.
- e.g. $\text{Poly}(\mathbb{R}^m, \mathbb{R}^n) \simeq \text{Lin}(\mathbb{R}[X_1, \dots, X_m], \mathbb{R}^n) \simeq \text{Lin}(\text{Sym}(\mathbb{R}^m), \mathbb{R}^n)$
- in categorical models: $!$ is a comonad on a symmetric monoidal category, satisfying some conditions
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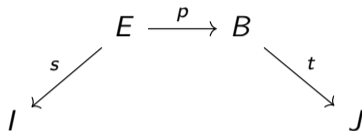
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Polynomials in categories

A polynomial from I to J in a category \mathcal{C} is a diagram



When $\mathcal{C} = \text{Set}$, it induces a **polynomial functor**

$$\begin{aligned} \text{Set}^I &\rightarrow \text{Set}^J \\ (X_i)_{i \in I} &\mapsto \left(\sum_{b \in t^{-1}(j)} \prod_{e \in p^{-1}(b)} X_{s(e)} \right)_{j \in J} \end{aligned}$$

- “ $B =$ monomials”
- “ $E =$ exponents/arities”

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$$\begin{array}{ccc} & E & \xrightarrow{p} & B \\ & \swarrow s & & \searrow t \\ I & & & J \end{array}$$

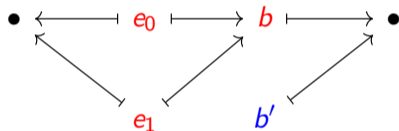
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Example

$$\{\bullet\} \xleftarrow{s} E \xrightarrow{p} B \xrightarrow{t} \{\bullet\}$$



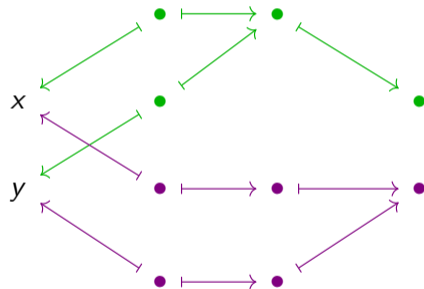
Induced functor:

Set \rightarrow Set

$$X \mapsto X^2 + 1$$

A more complicated example

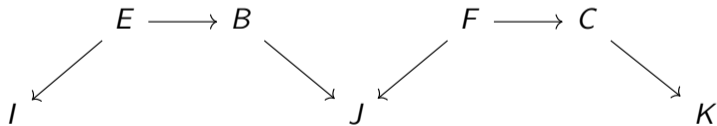
$$\{0, 1\} \xleftarrow{s} E \xrightarrow{p} B \xrightarrow{t} \{0, 1\}$$



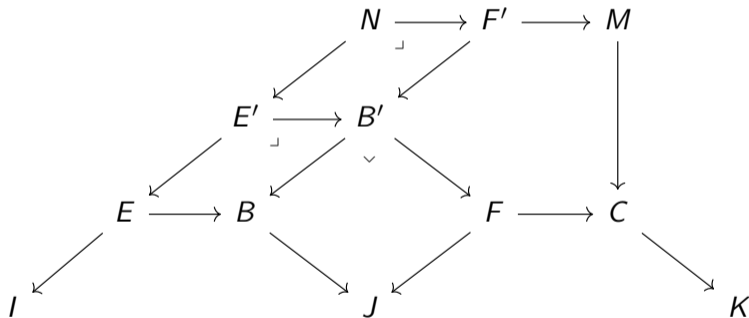
$$\text{Set}^2 \rightarrow \text{Set}^2$$

$$(X, Y) \mapsto (XY, X + Y)$$

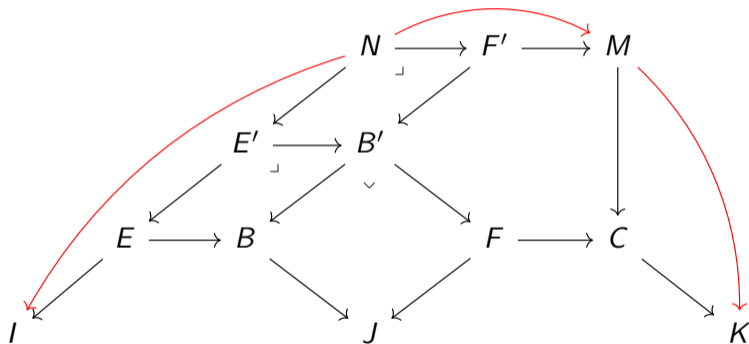
Composition of polynomials



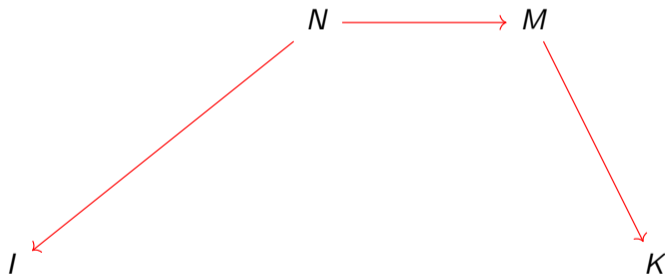
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Linear polynomials: spans

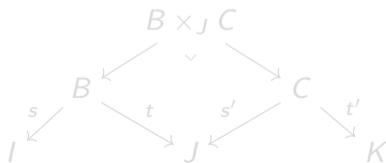
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is linear when p is an isomorphism: **the products are taken over singletons**

Linear polynomials are isomorphic to spans



And they compose via pullbacks



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$$\begin{array}{ccc} & B & \\ & \swarrow s & \searrow t \\ I & & J \end{array} \rightsquigarrow \begin{array}{ccc} & B \xlongequal{\text{id}_B} B & \\ & \swarrow s & \searrow t \\ I & & J \end{array}$$

And they compose via pullbacks

$$\begin{array}{ccccc} & & B \times_J C & & \\ & \swarrow & \downarrow & \searrow & \\ & B & & C & \\ & \swarrow s & & \swarrow s' & \searrow t' \\ I & & J & & K \end{array}$$

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The diagram shows two triangular diagrams connected by a wavy arrow. The left triangle has vertices I , B , and J with arrows $s: B \rightarrow I$ and $t: B \rightarrow J$. The right triangle has vertices I , $B \xrightarrow{\text{id}_B} B$, and J with arrows $s: B \rightarrow I$ and $t: B \rightarrow J$.

And they compose via pullbacks

The diagram shows a pullback square. At the top is the pullback object $B \times_J C$. Below it are two objects B and C . Arrows point from $B \times_J C$ to B and C . Below B and C are objects I , J , and K . Arrows point from B to I (labeled s) and J (labeled t). Arrows point from C to J (labeled s') and K (labeled t'). A vertical arrow points from $B \times_J C$ down to J .

The bicategory of polynomials

- Composition of polynomials uses pullbacks and other universal properties
- So it is not strictly associative/unital
- Polynomials thus form a **bicategory**
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- $\text{Poly}(\mathbb{R}^m, \mathbb{R}^n) \simeq \text{Lin}(\mathbb{R}[X_1, \dots, X_m], \mathbb{R}^n)$
- A basis of $\mathbb{R}[X_1, \dots, X_m]$ is given by monomials.
- In our categorified setting, we have arbitrary sets as exponents.
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Thus we restrict ourselves to **finitary polynomials**.

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Finitary polynomials

$$I \xleftarrow{s} E \xrightarrow{p} B \xrightarrow{t} J$$

is said to be **finitary** if $\forall b \in B, p^{-1}(b)$ is finite.

Examples:

- $X \mapsto X^3 + X + 1$ is finitary
- $X \mapsto \mathbb{N} \times X$ is finitary
- $(X_i)_{i \in \mathbb{N}} \mapsto ((X_i)^i)_{i \in \mathbb{N}}$ is finitary
- $X \mapsto X^{\mathbb{N}}$ is **not** finitary
- a linear polynomial is always finitary

(more generally, we could restrict to polynomials with arities in a fixed universe \mathcal{V} of “small” sets/types)

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Spans and polynomials

Span and Poly the bicategories of spans and finitary polynomials in sets.

- We would like $\text{Poly}(I, J) \simeq \text{Span}(!I, J)$
- Idea: monomials in $\mathbb{R}[X_1, \dots, X_m]$ are given by **multisets** over $\{1, \dots, m\}$
- Can we hope for $\text{Poly}(I, J) \simeq \text{Span}(\text{Mul}(I), J)$?

Yes and no... At the level of sets, yes, but Poly and Span are bicategories, and those groupoids of morphisms are not equivalent !

We need to :

- replace multisets by **homotopy multisets**,
- replace sets by groupoids.

To do that, we work in Homotopy Type Theory.

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- HoTT is an extension of Martin L of Type Theory where types are thought of as **spaces**.
- Spaces in the sense of **homotopy theory**.
- Discrete types are sets.
- Groupoids also are types, as are 2-groupoids, n -groupoids, ∞ -groupoids.
- In this context, Σ -types look like a Grothendieck construction.
- Quotients are **homotopy quotients** : instead of identifying elements, they paths between them.

$$\text{Bool} \begin{array}{c} \xrightarrow{\text{id}} \\ \xrightarrow{\neg} \end{array} \text{Bool} \longrightarrow \{\bullet\} \quad \text{coequalizer}$$

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- Groupoids also are types, as are 2-groupoids, n -groupoids, ∞ -groupoids.
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Spans and polynomials in HoTT

- Ordinary multisets: $\text{Mul}(X) := \sum_{n:\mathbb{N}} X^n / \Sigma_n$
- Equivalently, the **free commutative monoid on X** .
- Homotopy multisets: $\text{HMul}(X) := \sum_{n:\mathbb{N}} X^n // \Sigma_n$ (homotopy quotient).
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With this last definition, and using spans and polynomials in **types**, we proved in **HoTT**:

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A higher categorical model

- Switching from sets to groupoids makes Poly and Span into 3-categories.
- Going to arbitrary types, we get an ∞ -category : associativity and unitality up to isomorphisms, themselves satisfying coherence laws, etc.
- We cannot state or prove those infinite **homotopy coherence laws** in HoTT, so we work with **wild categories**.
- Wild categories have the standard definition of categories, but with sets replaced by types.
- No pentagon or triangle isomorphisms required of the associators and unitors.

Remark

Not all coherences can be stated in HoTT, but some can be proven meta-theoretically. For instance, we can prove a wild category has cartesian products, and know meta-theoretically that the induced monoidal structure is homotopy coherent.

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A full model of classical linear logic

- Poly is the Kleisli category of a comonad \mathbf{HMul} on \mathbf{Span} .
- \mathbf{HMul} makes \mathbf{Span} into a Seely category: a model of intuitionistic linear logic.
- \mathbf{Span} is moreover compact closed, with self-dual objects.
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Examples of higher polynomials - 1

$$\{\bullet\} \longleftarrow \mathbb{N} \xrightarrow{p} \sum_{n:\mathbb{N}} B(\mathbb{Z}/n\mathbb{Z}) \longrightarrow \{\bullet\}$$

$$n \longmapsto \{1, \dots, n\}$$

- $B(\mathbb{Z}/n\mathbb{Z})$ is the groupoid with one point and $\mathbb{Z}/n\mathbb{Z}$ as automorphisms
- $p^{-1}(\{1, \dots, n\}) \simeq \mathbb{Z}/n\mathbb{Z}$
- p^{-1} is taken in the sense of **homotopy fiber**
- Induced polynomial : $F(X) = \sum_{n:\mathbb{N}} X^n // (\mathbb{Z}/n\mathbb{Z})$
- The type of cyclic lists over X
- Generally, summing over groupoids amounts to quotienting the summand

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$$(E, e) \longmapsto E$$

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Examples of higher polynomials - The Hopf Fibration

$$\{\bullet\} \longleftarrow S^3 \xrightarrow{h} S^2 \longrightarrow \{\bullet\}$$

- The map H has fiber $h(x)$, merely equivalent to S^1 , the circle.
- $F(X) = \sum_{x:S^2} X^{h(x)}$.
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- If you have any idea what this represents, please reach out !

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We arranged the usual notions of spans and polynomials into a model of linear logic, using ideas from homotopy type theory. What next?

- Differential structure?
- Exploring other “homotopifications” of vector spaces and polynomials: spectra? stable ∞ -categories?
- Comparison with other span-based models of linear logic by Mellies, Clairambault, Forest
- Comparison with generalized species of structure